Propositional Logic 1: Motivation – Parse Tree

Mathematical Logic - First Term 2023-2024

MZI

School of Computing Telkom University

SoC Tel-U

September 2023

1 / 45

MZI (SoC Tel-U) Propositional Logic 1 September 2023

Acknowledgements

This slide is compiled using the materials in the following sources:

- Discrete Mathematics and Its Applications (Chapter 1), 8th Edition, 2019, by
 K. H. Rosen (primary reference).
- Discrete Mathematics with Applications (Chapter 2), 5th Edition, 2018, by S. S. Epp.
- Logic in Computer Science: Modelling and Reasoning about Systems (Chapter 1), 2nd Edition, 2004, by M. Huth and M. Ryan.
- Mathematical Logic for Computer Science (Chapter 2, 3, 4), 2nd Edition, 2000, by M. Ben-Ari.
- Discrete Mathematics 1 (2012) slides in Fasilkom UI by B. H. Widjaja.
- Mathematical Logic slides in Telkom University by A. Rakhmatsyah and B. Purnama.

Some figures are excerpted from those sources. This slide is intended for internal academic purpose in SoC Telkom University. No slides are ever free from error nor incapable of being improved. Please convey your comments and corrections (if any) to pleasedontspam>@telkomuniversity.ac.id.

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- Proposition: Definition
- Proposition: Example
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System Specification Consistency Problem

A software engineer is inquired by his manager to develop an information system that complies following specification:



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Can an information system with these specification be built?



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Can an information system with these specification be built? In other words, are these system specification consistent?

System specification consistency problem is one of the problems that can be solved using propositional logic which we are going to learn.

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Some Examples of Propositions

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$$3^4 - 4^3 < 10$$

- Is this a statement? Yes.
- Is this a proposition? Yes.
- Truth value? False (because $3^4 4^3 = 17$).



$$x + 3 \ge 2045$$



10 / 45

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- Is this a statement? Yes.
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10 / 45

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10 / 45

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10 / 45

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11 / 45

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Exercise



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Exercise

Verify whether following sentences are propositions or not.

"Have you understood the definition of a proposition?"



"Please read your textbooks regularly!"

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- "Have you understood the definition of a proposition?"
- (a) "I have understood the definition of a proposition."



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- "Have you understood the definition of a proposition?"
- "I have understood the definition of a proposition."
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- "Have you understood the definition of a proposition?"
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- "Bla bla bla, \$#@&%!"



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13 / 45

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- **1** unary operator; only requires one operand: negation $(\neg \text{ or } \sim)$;
- **@** binary operator; requires two operands: conjunction (\land) , disjunction (\lor) , exclusive disjunction/ exclusive-or (\oplus) , implication (\rightarrow) , bi-implication (\leftrightarrow)

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14 / 45

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The truth table for negation of a proposition:

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Exercise

Write in English the negation of following propositions:

- "I am a student"
- (2) "This month is not August"
- "Alex didn't do nothing"
- $2^{10} < 10^2$
- $3^4 \ge 4^3$

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15 / 45

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Example

Write in English the negation of following propositions.

- Bill is richer than Steve.
- Steve is older than Bill.

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- "It is not true that Steve is older than Bill" or in other words "Steve is as old as Bill or he is younger than Bill".

In Standard English, two negatives are understood to resolve to a positive; however, today as African American culture becomes more dispersed, the use of double negatives becomes increasingly prevalent and in many cases lead to confusions.

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• I didn't do nothing. (So, the speaker did something or not?)

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- I didn't do nothing. (So, the speaker did something or not?)
- ② I ain't got no money.

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For examples:

- I didn't do nothing. (So, the speaker did something or not?)
- ② I ain't got no money. (So, the speaker got some money or not?)

17 / 45

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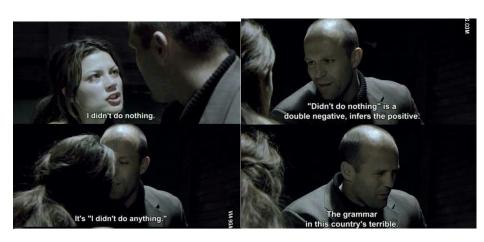
For examples:

- I didn't do nothing. (So, the speaker did something or not?)
- I ain't got no money. (So, the speaker got some money or not?)
- I'm not hungry no more.

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- 1 didn't do nothing. (So, the speaker did something or not?)
- I ain't got no money. (So, the speaker got some money or not?)
- I'm not hungry no more. (So, the speaker is hungry or not?)



Conjunction

Let p and q be propositions. Then $p \wedge q$ is also a proposition which is called the conjunction of p and q.

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19 / 45

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\mathbf{F}	\mathbf{T}	\mathbf{F}
F	F	F

Examples of Conjunctions

Exercise

Suppose we have following atomic propositions:

p: The sun rises in the east $q: 2 \times 3 \le 3^2$ r: Cat is a reptile $s: 2^4 > 4^2$

r: Cat is a reptile

Write and determine the truth value of the following compound propositions: (1) $p \wedge \neg q$; (2) $\neg r \wedge \neg s$.

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- ② $\neg r \wedge \neg s$: Cat is not a reptile and $2^4 \le 4^2$ Since "cat is not a reptile" is true and $2^4 \le 4^2$ is

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Propositional Logic 1

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Let p and q be propositions. Then $p \vee q$ is also a proposition which is called the disjunction of p and q.

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21 / 45

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Observe that $p \vee q$ is also true if both p and q are true.

MZI (SoC Tel-U) Propositional Logic 1 September 2023 21 / 45

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22 / 45

Exclusive Disjunction (XOR)

Exclusive Disjunction (*Exclusive-OR – XOR*)

Let p and q be propositions. Then $p \oplus q$ is also a proposition which is called the exclusive disjunction (exclusive or, xor) of p and q.

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Exercise

Suppose we have following atomic propositions:

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p: Alex likes cookies q: Alex likes pizza
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r:2 is an even number s:2 is a prime number

Write the compound proposition: (1) $p \oplus q$; (2) $r \oplus s$ and also determine the truth value of $r \oplus s$.

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Suppose we have following atomic propositions:

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Write the compound proposition: (1) $p \oplus q$; (2) $r \oplus s$ and also determine the truth value of $r \oplus s$.

Solution:

 \bullet $p \oplus q$: "Alex likes cookies but he doesn't like pizza, or Alex doesn't like cookies but he likes pizza";

24 / 45

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Suppose we have following atomic propositions:

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Solution:

• $p \oplus q$: "Alex likes cookies but he doesn't like pizza, or Alex doesn't like cookies but he likes pizza"; or we can also write "Alex likes either cookies or pizza, but not both".

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- ② $r \oplus s$: "2 is an even number but it is not a prime number, or 2 is not an even number but it is a prime number";

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Write the compound proposition: (1) $p \oplus q$; (2) $r \oplus s$ and also determine the truth value of $r \oplus s$.

- $p \oplus q$: "Alex likes cookies but he doesn't like pizza, or Alex doesn't like cookies but he likes pizza"; or we can also write "Alex likes either cookies or pizza, but not both".
- $oldsymbol{0}$ $r\oplus s$: "2 is an even number but it is not a prime number, or 2 is not an even number but it is a prime number"; or we can also write "2 is either an even number or a prime number, but not both". The truth of each of r and s

Exercise

Suppose we have following atomic propositions:

```
p: Alex likes cookies q: Alex likes pizza r: 2 is an even number s: 2 is a prime number
```

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Solution:

- $p \oplus q$: "Alex likes cookies but he doesn't like pizza, or Alex doesn't like cookies but he likes pizza"; or we can also write "Alex likes either cookies or pizza, but not both".
- ② $r \oplus s$: "2 is an even number but it is not a prime number, or 2 is not an even number but it is a prime number"; or we can also write "2 is either an even number or a prime number, but not both". The truth of each of r and s is true, therefore the truth of $r \oplus s$ is

24 / 45

Exercise

Suppose we have following atomic propositions:

```
p: Alex likes cookies q: Alex likes pizza r: 2 is an even number s: 2 is a prime number
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Write the compound proposition: (1) $p \oplus q$; (2) $r \oplus s$ and also determine the truth value of $r \oplus s$.

- $p \oplus q$: "Alex likes cookies but he doesn't like pizza, or Alex doesn't like cookies but he likes pizza"; or we can also write "Alex likes either cookies or pizza, but not both".
- ② $r \oplus s$: "2 is an even number but it is not a prime number, or 2 is not an even number but it is a prime number"; or we can also write "2 is either an even number or a prime number, but not both". The truth of each of r and s is true, therefore the truth of $r \oplus s$ is false.

Implication or Conditional Statement

Let p and q be propositions, then $p \to q$ is also a proposition which is called an implication or a conditional statement of the proposition "if p, then q". In an implication $p \to q$, p is called the hypothesis/ antecedent/ premise and q is called the conclusion/ consequence.

 $p \rightarrow q$ is read:

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p \rightarrow q is read:
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\begin{array}{c|cccc} \text{if } p \text{ then } q & & q \text{ if } p \\ p \text{ implies } q & & q \text{ is implied by } p \end{array}
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$p \rightarrow q$ is read:

```
\begin{array}{c|cccc} \text{if } p \text{ then } q & & q \text{ if } p \\ p \text{ implies } q & & q \text{ is implied by } p \\ p \text{ is sufficient for } q & q \text{ is necessary for } p \end{array}
```

$$\begin{array}{c|cc} p & q & p \to q \\ \hline T & T & \end{array}$$

$$\begin{array}{c|cc} p & q & p \to q \\ \hline T & T & T \\ T & F & \end{array}$$

$$\begin{array}{c|ccc} p & q & p \rightarrow q \\ \hline T & T & T \\ \hline T & F & F \\ \hline F & T & \end{array}$$

p	q	$p \rightarrow q$
T	T	${f T}$
\mathbf{T}	\mathbf{F}	\mathbf{F}
\mathbf{F}	\mathbf{T}	${f T}$
F	\mathbf{F}	

p	q	$p \rightarrow q$
\mathbf{T}	Τ	${f T}$
\mathbf{T}	\mathbf{F}	\mathbf{F}
\mathbf{F}	\mathbf{T}	${f T}$
F	\mathbf{F}	${f T}$

Example

Consider following propositions:

p: "Alex's final exam is 100"

q: "Alex's final grade is A"

 $p \rightarrow q$:

Example

Consider following propositions:

p: "Alex's final exam is 100"

q: "Alex's final grade is A"

 $p \rightarrow q$: "if Alex's final exam is 100, then

Example

Consider following propositions:

p: "Alex's final exam is 100"

q: "Alex's final grade is A"

 $p \rightarrow q$: "if Alex's final exam is 100, then his final grade is A"

 $p \rightarrow q$ is false (F) when

27 / 45

Example

Consider following propositions:

p: "Alex's final exam is 100"

q: "Alex's final grade is A"

 $p \rightarrow q$: "if Alex's final exam is 100, then his final grade is A"

 $p \rightarrow q$ is false (F) when Alex's final exam is 100,

Example

Consider following propositions:

p: "Alex's final exam is 100"

q: "Alex's final grade is A"

 $p \rightarrow q$: "if Alex's final exam is 100, then his final grade is A"

 $p \to q$ is false (F) when Alex's final exam is 100, but his final grade is not A"

 $p \rightarrow q$ is true (T) when:

Example

Consider following propositions:

p: "Alex's final exam is 100"

q: "Alex's final grade is A"

 $p \rightarrow q$: "if Alex's final exam is 100, then his final grade is A"

 $p \to q$ is false (F) when Alex's final exam is 100, but his final grade is not A"

 $p \rightarrow q$ is true (T) when:

Alex's final exam is not 100

Example

Consider following propositions:

p: "Alex's final exam is 100"

q: "Alex's final grade is A"

 $p \rightarrow q$: "if Alex's final exam is 100, then his final grade is A"

 $p \to q$ is false (F) when Alex's final exam is 100, but his final grade is not A"

 $p \rightarrow q$ is true (T) when:

- Alex's final exam is not 100
- Alex's final grade is A.

Exercise

At the beginning of a season, an owner of a basketball club once said, "If my club win this season's final, then I will sell this club". When the season ended, the club was relegated to the lower division, yet the owner still sold his club. Does the act of selling the club violate the owner's previous statement?

Contrapositive, Converse, and Inverse

Let $p \rightarrow q$ be an implication, then:

- ullet the **contrapositive** (sometimes called contraposition) of p o q is $\neg q o \neg p$
- ullet the **converse** of p o q is q o p
- the **inverse** of $p \rightarrow q$ is $\neg p \rightarrow \neg q$

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
\overline{T}	F	Т	F				

30 / 45

						contrapositive	converse	inverse
_	p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
_	T	F	T	F	${ m T}$			

30 / 45

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
Т	F	T	F	${ m T}$	T		

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	\mathbf{F}	T	T	T	

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
Т	F	Т	F	${ m T}$	T	T	T
\mathbf{T}	\mathbf{F}	F	\mathbf{T}		·		!

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
Т	F	Т	F	${ m T}$	T	T	T
\mathbf{T}	\mathbf{F}	F	${ m T}$	\mathbf{F}	'	1	ı

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	F	${ m T}$	T	T	T
\mathbf{T}	\mathbf{F}	\mathbf{F}	\mathbf{T}	\mathbf{F}	\mathbf{F}		'

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
Т	F	T	F	${ m T}$	T	T	T
\mathbf{T}	\mathbf{F}	\mathbf{F}	\mathbf{T}	\mathbf{F}	\mathbf{F}	${f T}$	

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	\mathbf{F}	T	T	T	T
\mathbf{T}	\mathbf{F}	F	\mathbf{T}	\mathbf{F}	\mathbf{F}	${f T}$	\mathbf{T}
\mathbf{F}	\mathbf{T}	T	\mathbf{F}		'		ı

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
Т	F	T	F	T	T	T	T
\mathbf{T}	\mathbf{F}	F	\mathbf{T}	\mathbf{F}	\mathbf{F}	${f T}$	\mathbf{T}
\mathbf{F}	${ m T}$	\mathbf{T}	\mathbf{F}	${f T}$	'	I	ı

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	F	T	T	T	T
\mathbf{T}	\mathbf{F}	\mathbf{F}	\mathbf{T}	\mathbf{F}	\mathbf{F}	${f T}$	\mathbf{T}
\mathbf{F}	\mathbf{T}	\mathbf{T}	\mathbf{F}	\mathbf{T}	${f T}$		'

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	F	${ m T}$	T	T	T
\mathbf{T}	\mathbf{F}	\mathbf{F}	\mathbf{T}	\mathbf{F}	\mathbf{F}	${f T}$	\mathbf{T}
\mathbf{F}	\mathbf{T}	\mathbf{T}	\mathbf{F}	${f T}$	${f T}$	\mathbf{F}	

						contrapositive	converse	inverse
	p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
_	T	F	T	\mathbf{F}	T	T	T	T
	\mathbf{T}	\mathbf{F}	\mathbf{F}	\mathbf{T}	F	\mathbf{F}	${f T}$	\mathbf{T}
	\mathbf{F}	\mathbf{T}	\mathbf{T}	\mathbf{F}	${f T}$	${f T}$	\mathbf{F}	\mathbf{F}
	\mathbf{F}	${f T}$	\mathbf{F}	${ m T}$		'	ı	I

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	\mathbf{F}	T	T	T	T
\mathbf{T}	\mathbf{F}	F	\mathbf{T}	F	F	${f T}$	\mathbf{T}
\mathbf{F}	\mathbf{T}	\mathbf{T}	\mathbf{F}	${f T}$	${f T}$	\mathbf{F}	\mathbf{F}
\mathbf{F}	\mathbf{T}	\mathbf{F}	${ m T}$	${f T}$			ı

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	F	T	T	T	T
\mathbf{T}	\mathbf{F}	\mathbf{F}	\mathbf{T}	\mathbf{F}	\mathbf{F}	${f T}$	\mathbf{T}
\mathbf{F}	\mathbf{T}	\mathbf{T}	\mathbf{F}	${f T}$	${f T}$	\mathbf{F}	\mathbf{F}
\mathbf{F}	\mathbf{T}	F	\mathbf{T}	\mathbf{T}	${f T}$!

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	\mathbf{F}	T	T	T	T
\mathbf{T}	\mathbf{F}	\mathbf{F}	\mathbf{T}	\mathbf{F}	\mathbf{F}	${f T}$	${f T}$
\mathbf{F}	\mathbf{T}	\mathbf{T}	\mathbf{F}	${f T}$	${f T}$	\mathbf{F}	\mathbf{F}
\mathbf{F}	\mathbf{T}	F	\mathbf{T}	\mathbf{T}	${f T}$	${ m T}$	

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	\mathbf{F}	T	T	T	T
\mathbf{T}	\mathbf{F}	F	\mathbf{T}	\mathbf{F}	\mathbf{F}	${f T}$	T
\mathbf{F}	\mathbf{T}	\mathbf{T}	\mathbf{F}	${f T}$	${f T}$	\mathbf{F}	\mathbf{F}
\mathbf{F}	\mathbf{T}	\mathbf{F}	\mathbf{T}	${f T}$	${f T}$	${f T}$	T

Observe that the truth of $p \rightarrow q$ is identical to the truth of

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	\mathbf{F}	T	T	T	T
\mathbf{T}	\mathbf{F}	F	${ m T}$	\mathbf{F}	\mathbf{F}	\mathbf{T}	\mathbf{T}
\mathbf{F}	\mathbf{T}	\mathbf{T}	\mathbf{F}	${f T}$	${f T}$	\mathbf{F}	\mathbf{F}
\mathbf{F}	\mathbf{T}	\mathbf{F}	\mathbf{T}	${f T}$	${f T}$	${f T}$	T

Observe that the truth of $p \to q$ is identical to the truth of $\neg q \to \neg p$, the truth of $q \to p$ is identical to the truth of

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	F	T	T	T	T
\mathbf{T}	\mathbf{F}	\mathbf{F}	\mathbf{T}	\mathbf{F}	\mathbf{F}	${f T}$	\mathbf{T}
\mathbf{F}	\mathbf{T}	\mathbf{T}	\mathbf{F}	${f T}$	${f T}$	\mathbf{F}	\mathbf{F}
\mathbf{F}	\mathbf{T}	F	\mathbf{T}	${f T}$	${f T}$	${f T}$	T

Observe that the truth of $p \to q$ is identical to the truth of $\neg q \to \neg p$, the truth of $q \to p$ is identical to the truth of $\neg p \to \neg q$. In this condition, we say

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	\mathbf{F}	T	T	T	T
\mathbf{T}	\mathbf{F}	F	\mathbf{T}	\mathbf{F}	\mathbf{F}	${f T}$	\mathbf{T}
\mathbf{F}	\mathbf{T}	T	\mathbf{F}	${f T}$	${f T}$	\mathbf{F}	\mathbf{F}
\mathbf{F}	\mathbf{T}	F	\mathbf{T}	\mathbf{T}	${f T}$	${f T}$	\mathbf{T}

Observe that the truth of $p\to q$ is identical to the truth of $\neg q\to \neg p$, the truth of $q\to p$ is identical to the truth of $\neg p\to \neg q$. In this condition, we say $p\to q$ is equivalent to $\neg q\to \neg p$ and

					contrapositive	converse	inverse
p	$\neg p$	q	$\neg q$	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$q \rightarrow p$	$\neg p \rightarrow \neg q$
T	F	T	\mathbf{F}	T	T	T	T
\mathbf{T}	\mathbf{F}	F	\mathbf{T}	\mathbf{F}	\mathbf{F}	${f T}$	\mathbf{T}
\mathbf{F}	\mathbf{T}	T	\mathbf{F}	${f T}$	${f T}$	\mathbf{F}	\mathbf{F}
\mathbf{F}	\mathbf{T}	F	\mathbf{T}	\mathbf{T}	${f T}$	${f T}$	\mathbf{T}

Observe that the truth of $p \to q$ is identical to the truth of $\neg q \to \neg p$, the truth of $q \to p$ is identical to the truth of $\neg p \to \neg q$. In this condition, we say $p \to q$ is equivalent to $\neg q \to \neg p$ and $q \to p$ is equivalent to $\neg p \to \neg q$.

Remark

Every conditional statement is equivalent to its contrapositive.

Bi-implication

Let p and q be propositions, then $p \leftrightarrow q$ is also a proposition which is called a bi-implication or a biconditional statement of the propositions "p if and only if q".

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p if and only if q if p then q, and conversely p is equivalent to q
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p iff q, (iff is a shorthand for if and only if) p is necessary and sufficient for q p and q are equivalent
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 $p \leftrightarrow q$ is true (T) precisely when $p \to q$ and $q \to p$ are both true (T).

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$$\begin{array}{c|cc} p & q & p \leftrightarrow q \\ \hline T & T & \end{array}$$

 $p \leftrightarrow q$ is true (T) precisely when $p \to q$ and $q \to p$ are both true (T).

$$\begin{array}{c|cc} p & q & p \leftrightarrow q \\ \hline T & T & T \\ T & F & \end{array}$$

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$$\begin{array}{c|ccc} p & q & p \leftrightarrow q \\ \hline T & T & T \\ T & F & F \\ F & T & \end{array}$$

 $p \leftrightarrow q$ is true (T) precisely when $p \to q$ and $q \to p$ are both true (T).

p	q	$p \leftrightarrow q$
T	T	${ m T}$
\mathbf{T}	\mathbf{F}	\mathbf{F}
\mathbf{F}	\mathbf{T}	\mathbf{F}
\mathbf{F}	\mathbf{F}	

 $p \leftrightarrow q$ is true (T) precisely when $p \to q$ and $q \to p$ are both true (T).

Truth table for a biconditional statement:

p	q	$p \leftrightarrow q$
T	T	T
\mathbf{T}	\mathbf{F}	\mathbf{F}
\mathbf{F}	\mathbf{T}	\mathbf{F}
\mathbf{F}	\mathbf{F}	${f T}$

Example

Consider following propositions:

p: "Alex's final grade was not less than 50"

q: "Alex passed the class"

 $p \leftrightarrow q$:

Example

Consider following propositions:

p: "Alex's final grade was not less than 50"

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 $p \leftrightarrow q$: "Alex's final grade was not less than 50 if and only if he passed the class"

 $p \leftrightarrow q$

Example

Consider following propositions:

p: "Alex's final grade was not less than 50"

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 $p \leftrightarrow q$ is true (T) when

 \bullet Alex's final grade was not less than 50 and he passed the class; or

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 $p \leftrightarrow q$: "Alex's final grade was not less than 50 if and only if he passed the class"

 $p \leftrightarrow q$ is true (T) when

- lacktriangle Alex's final grade was not less than 50 and he passed the class; or
- Alex's final grade

Example

Consider following propositions:

p: "Alex's final grade was not less than 50"

q: "Alex passed the class"

 $p \leftrightarrow q$: "Alex's final grade was not less than 50 if and only if he passed the class"

 $p \leftrightarrow q$ is true (T) when

- $oldsymbol{0}$ Alex's final grade was not less than 50 and he passed the class; or
- $oldsymbol{0}$ Alex's final grade was less than 50

Example

Consider following propositions:

p: "Alex's final grade was not less than 50"

q: "Alex passed the class"

 $p \leftrightarrow q$: "Alex's final grade was not less than 50 if and only if he passed the class"

 $p \leftrightarrow q$ is true (T) when

- $oldsymbol{0}$ Alex's final grade was not less than 50 and he passed the class; or
- ② Alex's final grade was less than 50 and he didn't pass the class.

 $p \leftrightarrow q$ is false (F) when

Example

Consider following propositions:

p: "Alex's final grade was not less than 50"

q: "Alex passed the class"

 $p \leftrightarrow q$: "Alex's final grade was not less than 50 if and only if he passed the class"

 $p \leftrightarrow q$ is true (T) when

- $oldsymbol{0}$ Alex's final grade was not less than 50 and he passed the class; or
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 $p \leftrightarrow q$ is true (T) when

- lacktriangle Alex's final grade was not less than 50 and he passed the class; or
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Example

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 $p \leftrightarrow q$: "Alex's final grade was not less than 50 if and only if he passed the class"

 $p \leftrightarrow q$ is true (T) when

- f O Alex's final grade was not less than 50 and he passed the class; or
- ② Alex's final grade was less than 50 and he didn't pass the class.

 $p \leftrightarrow q$ is false (F) when

Alex's final grade was not less than 50, but he didn't pass the class; or

Example

Consider following propositions:

- p: "Alex's final grade was not less than 50"
- q: "Alex passed the class"
- $p \leftrightarrow q$: "Alex's final grade was not less than 50 if and only if he passed the class"
- $p \leftrightarrow q$ is true (T) when
 - lacktriangle Alex's final grade was not less than 50 and he passed the class; or
 - $oldsymbol{0}$ Alex's final grade was less than 50 and he didn't pass the class.
- $p \leftrightarrow q$ is false (F) when
 - Alex's final grade was not less than 50, but he didn't pass the class; or
 - Alex's final grade

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 - lacktriangle Alex's final grade was not less than 50 and he passed the class; or
 - $oldsymbol{0}$ Alex's final grade was less than 50 and he didn't pass the class.
- $p \leftrightarrow q$ is false (F) when
 - Alex's final grade was not less than 50, but he didn't pass the class; or
 - ② Alex's final grade was less than 50,

Example

Consider following propositions:

- p: "Alex's final grade was not less than 50"
- q: "Alex passed the class"
- $p \leftrightarrow q$: "Alex's final grade was not less than 50 if and only if he passed the class"
- $p \leftrightarrow q$ is true (T) when
 - lacktriangle Alex's final grade was not less than 50 and he passed the class; or
 - $oldsymbol{0}$ Alex's final grade was less than 50 and he didn't pass the class.
- $p \leftrightarrow q$ is false (F) when
 - Alex's final grade was not less than 50, but he didn't pass the class; or
 - 4 Alex's final grade was less than 50, but he passed the class.

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- Motivation
- Proposition: Definition
- Proposition: Example
- 4 Logical Operators and Compound Propositions
- Precedences of Logical Operators
- 6 Propositional Formulas (Supplementary)

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35 / 45

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Which one of the following propositions is equivalent to $p \wedge q \rightarrow r$?

- $(p \land q) \to r$



The precedence of logical operators determines which operator has to be operated first (i.e., be applied to an operand)

36 / 45

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The precedence of logical operators is described by following table:

Operator	Precedence
	1
\wedge	2
\vee	3
\oplus	4
\rightarrow	5
\longleftrightarrow	6

36 / 45

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The precedence of logical operators is described by following table:

Operator	Precedence
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\oplus	4
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\longleftrightarrow	6

As in arithmetic expressions, we use parentheses "(" and ")" to specify or clarify the order of an operation in a compound proposition.

Put the parentheses to clarify the precedences of logical operators in the following compound propositions:

- $\bigcirc \neg p \lor q$

Solution:

 $\bigcirc p \lor q \land r \text{ means}$

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37 / 45

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Propositional Formulas

Propositional Formulas

A propositional formula may consist of:

- propositional constant: T (true) or F (false)
- atomic variables:

$$p, p_1, p_2, \dots$$

 q, q_1, q_2, \dots
 r, r_1, r_2, \dots

- **o** propositional operators: \neg , \land , \lor , \oplus , \rightarrow , \leftrightarrow
- and is constructed using following rules:
 - every atomic proposition is a propositional formula
 - ② if A and B are propositional formulas, then $\neg A$, $A \wedge B$, $A \vee B$, $A \oplus B$, $A \to B$, $A \leftrightarrow B$, are all propositional formulas.

39 / 45

Example

- $lackbox{0}$ $p \wedge q$ is a propositional formula
- $oldsymbol{arrho}$ pqee is not a propositional formula

Example

- lacksquare $p \wedge q$ is a propositional formula
- 2 $pq\lor$ is not a propositional formula
- \bullet $\neg\neg\,(\neg p\to\neg\neg r)$ is a propositional formula that can be rewritten as $\neg\,(\neg\,(\neg p\to\neg\,(\neg r)))$

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- $lackbox{0}$ $p \wedge q$ is a propositional formula
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- $\bullet \ p \oplus p \vee q \to r \wedge s \text{ is a propositional formula that can be rewritten as } (p \oplus (p \vee q)) \to (r \wedge s)$
- $\bigcirc \neg (\neg (\neg p \rightarrow q) \rightarrow r) \rightarrow s)$ is a propositional formula



Example

According to the definition of propositional formulas, we infer that

- $lackbox{0}$ $p \wedge q$ is a propositional formula
- 2 $pq\lor$ is not a propositional formula
- \bullet $\neg\neg\,(\neg p\to\neg\neg r)$ is a propositional formula that can be rewritten as $\neg\,(\neg\,(\neg p\to\neg\,(\neg r)))$
- \bullet $p \land q \rightarrow \oplus r \lor s$ is not a propositional formula
- $\ \, \textbf{0} \ \, p \vee q \vee \rightarrow r \oplus s \,\, \text{is not a propositional formula}$
- $\bullet p \oplus p \lor q \to r \land s \text{ is a propositional formula that can be rewritten as } (p \oplus (p \lor q)) \to (r \land s)$
- $\bigcirc \neg (\neg (\neg p \rightarrow q) \rightarrow r) \rightarrow s)$ is a propositional formula

4 D > 4 D > 4 E > 4 E > E = 40 Q Q

Subformula

Subformula

- lacktriangle A formula A is a subformula of A itself.
- ② If A and B are two propositional formulas used to construct a more complex propositional formula C, then A and B are proper subformulas of C.
- **9** Subformula is transitive: if A is a subformula of B and B is a subformula of C, then A is a subformula of C.

Example

let A be a formula $(p \land q) \rightarrow (r \lor s)$, then the subformulas of A are: (1)

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let A be a formula $(p \land q) \to (r \lor s)$, then the subformulas of A are: (1) $(p \land q) \to (r \lor s)$, (2)



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Example

let A be a formula $(p \land q) \rightarrow (r \lor s)$, then the subformulas of A are: (1) $(p \land q) \rightarrow (r \lor s)$, (2) $p \land q$, (3) $r \lor s$, (4) p, (5) q, (6)

41 / 45

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List all subformulas of $(p \rightarrow q) \lor (q \rightarrow p)$

Solution:



List all subformulas of $(p \rightarrow q) \lor (q \rightarrow p)$

$$\bullet \ (p \to q) \lor (q \to p)$$

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- $\bullet \ (p \to q) \lor (q \to p)$
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- $p \to q$
- $q \to p$
- p

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- $\bullet \ (p \to q) \lor (q \to p)$
- $p \to q$
- $q \to p$
- p
- q

List all subformulas of $(\neg p \land q) \rightarrow (p \land (q \lor \neg r))$

Solution:



List all subformulas of $(\neg p \wedge q) \to (p \wedge (q \vee \neg r))$

$$\bullet \ (\neg p \land q) \to (p \land (q \lor \neg r))$$

List all subformulas of $(\neg p \land q) \to (p \land (q \lor \neg r))$

- $\bullet \ (\neg p \wedge q) \to (p \wedge (q \vee \neg r))$
- $\bullet \ (\neg p \land q)$

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- $\bullet \ (p \wedge (q \vee \neg r))$



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- $\bullet \ q \vee \neg r$
- ¬p

List all subformulas of $(\neg p \land q) \rightarrow (p \land (q \lor \neg r))$

- $\bullet \ (\neg p \land q) \to (p \land (q \lor \neg r))$
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- ¬p
- ¬r
- p
- q

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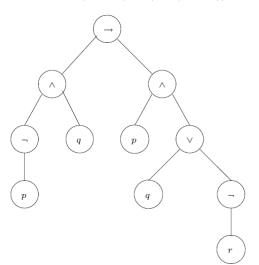
- $\bullet (\neg p \land q) \to (p \land (q \lor \neg r))$
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- $\bullet \ (p \land (q \lor \neg r))$
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- ¬p
- ¬r
- p
- .
- q
- \bullet r

Parse Tree

Parse tree can be used to visualize the structure of a propositional formula. For instance, the parse tree for $(\neg p \land q) \to (p \land (q \lor \neg r))$ is

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Draw the parse tree of each of these formulas: